Modular population protocols

Michael Raskin

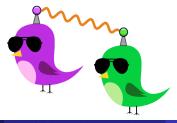
LaBRI, University of Bordeaux, CNRS UMR 5800

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Small sensors: constant memory, identical, pairwise interaction when accidentally near

Accepting and rejecting states Aiming for eventual consensus — unanimous acceptance or rejection (of initial configuration)

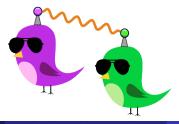
[AADFP 2004]



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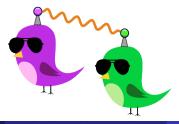
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Original motivation — small sensors, too small to be plausible today

Edge of decidability Even beefy servers like simplicity! Original motivation — small sensors, too small to be plausible today

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Eventual consensus on a simple predicate over unchanging inputs — limited ambitions

Also, most modular constructions are completely ad-hoc

Aims:

- Modularity
- Generality of tasks

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- Servers need to be assigned roles
- Roles have configuration requirements
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- Can schedule server for shutdown and maintenance/removal
- Can add servers
- Can change server configuration

Agents have inputs they can't change, and internal states Step options:

- Normal protocol step
- Reconfiguration

Reconfigurations:

Special input \perp means request to shutdown, special state \perp means the agent did shut down

- Add an agent that is shut down and has shutdown input
- Remove an agent that is shut down and has shutdown input
- Change an agent's input

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Two parts to specify

- Single-agent (local) input-output compatibility («32 GB storage not enough for storage servers»)
- Multiset-of-inputs to multiset-of-outputs (global) compatibility («Most of the servers have enough storage? Then webservers-to-storage-servers should be an equal split»)

After reconfigurations cease, each fair execution should eventually:

- Satisfy shutdown requests
- Stabilise all agent outputs
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Given two protocols implementing two specifications, can compose them:

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This implements composition of specifications: global and local compatibility relations are composed as multi-valued functions Given two protocols implementing two specifications, can compose them:

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Expressive power

A specification can be unimplementable for specific multisets of inputs

 $\ensuremath{\mathsf{wExact}}$ same number of web and storage servers, even if the total is odd»

Otherwise, if global compatibility is semilinear, the specification can be implemented by population protocols with reconfiguration

A specification may be implementable for wrong reasons «Only web front-ends, or a prime number of storage servers»

For any specification implemented by a protocol, another specification exists:

- Implemented by another protocol
- More restrictive (forbids everything forbidden by the original)
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- Cover more of the interesting tasks...
- ... and gain modularity
- Expressive power is still «semilinear»
- Future work: redefine and generalize known «efficient» protocols in this form

Thanks for your attention

Questions?

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