

# On Essentially Conditional Information Inequalities

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**Shannon type** inequalities [combinations of basic ineq]:

**example 1:**  $H(a) \leq H(a | b) + H(a | c) + I(b : c)$

**example 2:**  $2H(a, b, c) \leq H(a, b) + H(a, c) + H(b, c)$

## Linear information inequalities

General form: A linear information inequality is a combination of reals  $\{\lambda_{i_1, \dots, i_k}\}$  such that

$$\sum \lambda_{i_1, \dots, i_k} H(a_{i_1}, \dots, a_{i_k}) \geq 0$$

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Applications:

- multi-source network coding
- secret sharing
- combinatorial interpretations
- group theoretical interpretation
- Kolmogorov complexity
- ...

## Shannon type information inequalities:

- subadditivity  $H(A \cup B) \leq H(A) + H(B)$ ,
- submodularity  
$$H(A \cup B \cup C) + H(C) \leq H(A \cup C) + H(B \cup C),$$
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**Th** [Z. Zhang, R.W. Yeung 1998] There exists a non-Shannon type information inequality:

$$I(c : d) \leq 2I(c : d | a) + I(c : d | b) + I(a : b) \\ + I(a : c | d) + I(a : d | c)$$

**Theorem** [Z. Zhang, R.W. Yeung 1997] There exists a conditional non Shannon type inequality:

$$I(a : b) = I(a : b | c) = 0$$



$$I(c : d) \leq I(c : d | a) + I(c : d | b)$$

## Conditional information inequalities

(a) **Trivial, Shannon-type:**

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this is true since  $H(c) \leq H(c | a) + H(c | b) + I(a : b)$

[Shannon-type unconditional inequality]

## Conditional information inequalities

(b) **Trivial, non Shannon-type:**

if  $I(c : d | e) = I(e : c | d) = I(e : d | c) = 0$  then

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this is true since

$$\begin{aligned} I(c : d) &\leq I(c : d | a) + I(c : d | b) + I(a : b) \\ &\quad + I(c : d | e) + I(e : c | d) + I(e : d | c) \end{aligned}$$

[non Shannon-type unconditional inequality]

## Conditional information inequalities

(c) Non trivial, non Shannon-type:

- **Zhang, Yeung 97:** if  $I(a : b) = I(a : b | c) = 0$  then
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- **our result:** if  $H(c | a, b) = I(a : b | c) = 0$  then

$$I(c : d) \leq I(c : d | a) + I(c : d | b) + I(a : b)$$

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**Main Theorem.** These three statements are *essentially* conditional inequalities.

## Main Theorem [the first case of three]

The inequality

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More precisely, for all  $C_1, C_2$  the inequality

$$I(c : d) \leq I(c : d|a) + I(c : d|b) + C_1 I(a : b) + C_2 I(a : b|c)$$

does not hold!

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More precisely, for all  $C_1, C_2$  there exist  $(a, b, c, d)$  such that

$$I(c : d) \not\leq I(c : d|a) + I(c : d|b) + C_1 I(a : b) + C_2 I(a : b|c)$$

**Claim:** For any  $C_1, C_2$  there exist  $(a, b, c, d)$  such that

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Proof:

$a$	$b$	$c$	$d$	Prob[ $a, b, c, d$ ]
0	0	0	1	$(1 - \varepsilon)/4$
0	1	0	0	$(1 - \varepsilon)/4$
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$$\Theta(\varepsilon) \quad 0 \quad + \quad 0 \quad + \quad O(\varepsilon^2) \quad + \quad 0$$

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- Exactly the same classes of **unconditional** linear inequalities hold for Shannon's entropy and for Kolmogorov complexity.
- **Essentially conditional** inequality [ZY97] is true for Shannon's entropy **but not** for Kolmogorov complexity (in some natural sense).

See Proceedings for details.

**Question 1:** Can we apply essentially conditional inequalities (converse coding theorems, secret sharing problems, etc.)?

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**Question 2:** *May* we apply essentially conditional inequalities in 'real world' problems? These inequalities are not robust; do they make any 'physical' sense?

**Acknowledgments:** We indebted to the anonymous referees for helpful comments; following their suggestions we changed the title, reworked the introduction, modified statements of theorems, the bibliography, ..., and even corrected several *really essential* things in our paper.

Thank you! Any questions?